

Assignment 13
Published on 2020-12-06,
Estimated Time: 30 minutes,
Max.grade 10%

1 Dijkstra's Algorithm

(Goodrich2011, p.640) defines Dijkstra's algorithm. See also <https://bit.ly/2JSXqMU>. It is an efficient algorithm; it requires $O((m+n) \log_2 n)$ time, if we use priority queues; here m is the number of edges and n is the number of vertices in a graph.

In this exercise you do not need to implement a priority queue; assume that you can always pick the vertex with the smallest distance and add it to the set S of visited vertexes (those having distances already computed).

2 Problem

We start with the graph shown in Figure 1.

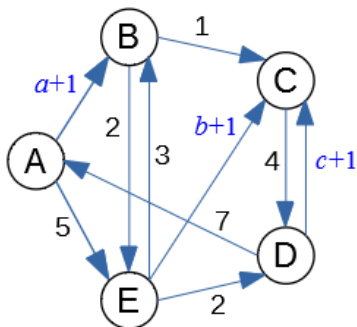


Figure 1: Graph diagram

The following edges (A, B) , (E, C) , (D, C) have weights $a + 1$, $b + 1$, and $c + 1$ respectively (here a, b, c should be replaced by the digits from your Student ID). Vertex A will be your source vertex. (You can assume that the distance from A to itself is 0; initially all the other distances are infinite, but then Dijkstra's algorithm relaxes them).

(A) Redraw the graph Figure 1, replace the edge weights $a + 1$, $b + 1$, and $c + 1$ by your values of a, b, c .

(B) Run the Dijkstra's algorithm; create a table showing how distances to A, B, C, D, E

change as the relaxations are performed. At every iteration highlight which vertex (among those not yet finished) has the minimum distance. Add it to the set S of finished vertices (the set S will have 0 in the very first iteration; after that it will grow by one vertex at a time).

(C) Summarize the result: For each of the 5 vertices tell what is its minimum distance from the source. Also tell what is the shortest path how to get there.